

## Conflict-Avoiding Codes of Length $n = 4m$ and Weight 3

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A *conflict-avoiding code* of length  $n$  for  $k$  active users is a set  $C \subseteq \{0, 1\}^n$  of binary vectors, called *codewords*, all of Hamming weight  $k$ , such that the Hamming distance between arbitrary cyclic shifts of distinct codewords is at least  $2k - 2$ . The support of a codeword is the set of indices of its coordinate positions with 1. A conflict-avoiding code of weight  $k$  is said to be *equi-difference* (when  $k = 3$ , it is said to be *centered*, or *symmetric*), if the supports of all codewords are of the form  $\{0, i, 2i, \dots, (k-1)i\}$  for some  $i \in \mathbb{Z}_n \setminus \{0\}$ . A conflict-avoiding code of the maximum size is said to be *optimal*.

In this talk, we will briefly glance at the technical background first, and then by using a kind of graph and linear programming, we will improve previously known upper bounds on the size of an optimal (centered) conflict-avoiding code of length  $n = 4m$  and weight 3. Furthermore, it will be proved that the improved upper bounds are sharp for some series of  $m$  by direct constructions.